# Staged generics-sop 

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2020-01-07

## AWell-Typed

The Haskell Consultants

## generics-sop

## Sums and products

Sum :: (a -> Type) -> [a] -> Type
Product :: (a -> Type) -> [a] -> Type

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Sum $f\left[x_{1}, x_{2}, x_{3}\right] \approx f x_{1}+f x_{2}+f x_{3}$ Product $f\left[x_{1}, x_{2}, x_{3}\right] \approx f x_{1} \times f x_{2} \times f x_{3}$

## Sums and products

Sum :: (a -> Type) -> [a] -> Type
Product :: (a -> Type) -> [a] -> Type

Sum $f\left[x_{1}, x_{2}, x_{3}\right] \approx f x_{1}+f x_{2}+f x_{3}$
Product $f\left[x_{1}, x_{2}, x_{3}\right] \approx f x_{1} \times f x_{2} \times f x_{3}$

Sum (Product f) ([ $\left.x_{1}, x_{2}\right],[],\left[x_{3}, x_{4}, x_{5}\right]$ )

$$
\approx\left(f x_{1} \times f x_{2}\right)+1+\left(f x_{3} \times f x_{4} \times f x_{5}\right)
$$

## Example datatype

data Animal =
HoppingAnimal String Double
| WalkingAnimal String Int

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Sum (Product I) [[String, Double], [String, Int]] $\approx$ (I String $\times$ I Double) + (I String $\times$ I Int) $\approx$ (String $\times$ Double) + (String $\times$ Int $)$

## Example datatype

data Animal =
HoppingAnimal String Double | WalkingAnimal String Int

Sum (Product I) [[String, Double], [String, Int]] $\approx$ (I String $\times$ I Double) + (I String $\times$ I Int) $\approx$ (String $\times$ Double) + (String $\times$ Int $)$

Description Animal = [[String, Double], [String, Int]]
from :: Animal -> Sum (Product I) (Description Animal)
to :: Sum (Product I) (Description Animal) -> Animal

## A class for representable types

```
class (All (All Top) (Description a)) => Generic a where
    type Description a :: [[Type]]
    from :: a -> Sum (Product I) (Description a)
    to :: Sum (Product I) (Description a) -> a
```


## Operations on sums and products

## mapsum : :

$$
\begin{aligned}
& \text { All Top xs } \\
\Rightarrow & (\forall x \cdot f x->g x)->\text { Sum } \quad f \times s->\text { Sum } \quad g \times s
\end{aligned}
$$

## mapproduct :

$$
\begin{aligned}
& \text { All Top } x s \\
\Rightarrow & (\forall x \cdot f x->g x)->\text { Product } f x s->\text { Product } g x s
\end{aligned}
$$

## Operations on sums and products

cmap $_{\text {Sum }} \quad:$ :

$$
\begin{aligned}
& \text { All c xs } \\
\Rightarrow & (\forall x \cdot c x=>f x \rightarrow g x) \\
\Rightarrow & \text { Sum } \quad f \times s \rightarrow \text { Sum } \quad g x s
\end{aligned}
$$

cmapproduct : :
All c xs
$\Rightarrow(\forall x . c x=>f x \rightarrow g x)$
-> Product f xs -> Product g xs
cmap $_{\text {SoP }} \quad::$

$$
\begin{aligned}
& \text { All (All c) xs } \\
& \Rightarrow(\forall x . c x=>f x \rightarrow g x) \\
& \text {-> Sum (Product f) xs -> Sum (Product g) xs }
\end{aligned}
$$

## Operations on sums and products

$$
\begin{aligned}
& \text { cpur } \text { Product }^{\text {Pro }} \\
& \text { All c xs } \\
& \quad=>(\forall x \cdot c x=>f x) \\
& \quad>\text { Product } f x s
\end{aligned}
$$

## Operations on sums and products

$$
\begin{aligned}
& \text { collapsesum : All Top xs => Sum (K a) xs -> a } \\
& \text { collapse }{ }_{\text {Product }} \text { : : All Top xs => Product (K a) xs -> [a] }
\end{aligned}
$$

## Operations on sums and products

zipWithproduct : :

$$
\begin{aligned}
& \text { All Top xs } \\
=> & (\forall x \cdot f \times x->g x->h x) \\
-> & \text { Product } f \times s \rightarrow \text { Product } g \text { xs }->\text { Product } h x s
\end{aligned}
$$

Operations on sums and products

```
zipWithProduct ::
        All Top xs
    => (}\forall\textrm{x}.\textrm{f}x -> g x -> h x)
    -> Product f xs -> Product g xs -> Product h xs
zipWithsum ::
        All Top xs
    => (}\forall\textrm{x}.f\textrm{f}|>\textrm{g}x>>hx
    -> Product f xs -> Sum g xs -> Sum h xs
```


## Operations on sums and products

```
anaproduct ::
        All Top xs
    => (\forally ys . s (y : ys) -> (f y, s ys))
    -> s xs -> Product f xs
```


## Arities of each constructor

```
constructorArities : :
    Generic a => Product (K Word) (Description a)
constructorArities =
    cpure Product \(^{( }\)(All Top) go
    where
    go : : \(\forall\) xs . All Top xs \(=>K\) Word \(x s\)
go \(=K(\) fromIntegral (lengthsList \(@ x s))\)
```


## Arities of each constructor

constructorArities : :
Generic a => Product (K Word) (Description a)
constructorArities =
cpureproduct @(All Top) go where

$$
\begin{aligned}
& \text { go }:: \forall \text { xs . All Top xs }=>\text { K Word xs } \\
& \text { go }=K(\text { fromIntegral (length SList @xs)) }
\end{aligned}
$$

Example:
data Animal =
HoppingAnimal String Double
| WalkingAnimal String Int
constructorArities @Animal

$$
=\text { K } 2: * \text { K } 2: * \mathrm{Nil}
$$

## Numbering each constructor

```
constructorNumbers ::
    Generic a => Product (K Word) (Description a)
constructorNumbers =
    anaproduct
        (\ (K i) -> (K i, K (i + 1)))
    (K 0)
```


## Numbering each constructor

constructorNumbers : : Generic a => Product (K Word) (Description a) constructorNumbers = anaproduct

```
(\ (K i) -> (K i,K (i + 1)))
(K 0)
```

Example:
data Animal =
HoppingAnimal String Double
| WalkingAnimal String Int
constructorNumbers @Animal

$$
=\mathrm{K} 0: * \mathrm{~K} 1: * \mathrm{Nil}
$$

## Encoding, generically

```
gencode ::
    \forall a . (Generic a, All (All Serialise) (Description a))
    => a -> Encoding
gencode x =
    collapsesum
    (czipWith3sum @(All Top)
        (\ (K a) (K i) encs ->
            K ( encodeListLen (a + 1)
            <> encodeWord i
            <> mconcat (collapse Product encs)
                )
            )
            (constructorArities @a)
            (constructorNumbers @a)
        (cmapsop @Serialise (\ (I y) -> K (encode y)) (from x))
    )
```


## Staging using Typed Template Haskell

## Quotes and splices

```
type Code a = Q (TExp a)
newtype Code' a = Code {unCode :: Code a}
```


## Quotes and splices

```
type Code a = Q (TExp a)
newtype Code' a = Code {unCode :: Code a}
ex 1 :: Code Int
ex = [|| 1 + 2 + 3 ||]
```


## Quotes and splices

```
type Code a = Q (TExp a)
newtype Code' a = Code {unCode :: Code a}
ex
ex}\mp@subsup{1}{1}{=[|| 1 + 2 + 3 ||]
ex2 :: Code Int
ex = [|| $$ex * * $$ex | ||]
AST:(1 + 2 + 3) * (1 + 2 + 3)
```


## Lifting

f : : Int -> Code Int
$f x=[||x+x||]$
$\mathrm{ex}_{3}$ :: Code Int
$e x_{3}=[\| \| \$(f(1+2+3)) \|]$
AST: $6+6$

## Lifting

```
f :: Int -> Code Int
f x = [|| x + x ||]
ex3 :: Code Int
ex3 = [|| $$(f (1 + 2 + 3)) ||]
AST: 6 + 6
g :: Lift a => [a] -> Code [a]
g xs = [|| reverse xs ||]
ex4 :: Code [Int]
ex4 = [|| $$(g (replicate 3 1)) ||]
AST: reverse [1, 1, 1]
```


## Using variables before they are defined

$$
\begin{aligned}
& f:: \text { Int -> Code Int } \\
& \text { f } x=[||x+x||] \\
& e x_{5}:: \text { Code (Int -> Int) } \\
& e x_{5}=[||\backslash x->\$ \$(f x)||] \text {-- not ok }
\end{aligned}
$$

Stage error: ' x ' is bound at stage 2 but used at stage 1

## Using variables before they are defined

```
f :: Int -> Code Int
f x = [|| x + x ||]
ex 
ex = [|| \x -> $$(f x) ||] -- not ok
```

Stage error: ' x ' is bound at stage 2 but used at stage 1

But this is ok:
h :: Code Int -> Code Int
h $x=[||\$ \$ x+\$ \$ 1|]$
ex 6 :: Code (Int -> Int)
$e x_{6}=[||\backslash x \rightarrow \$ \$(h[| | x| |])||]$
AST: $\backslash x->x+x$

## Hello world of staging

```
square : : Int -> Int square \(\mathrm{x}=\mathrm{x}\) * x
```


## Hello world of staging

```
square :: Int -> Int
square x = x * x
```

power : : Int -> Int -> Int
power n x
| $\mathrm{n}=0=1$
| even $n=$ square (power (n 'div' 2) $x$ )
| otherwise $=x$ * power $(n-1) x$

## Hello world of staging

```
square :: Int -> Int
square x = x * x
power :: Int -> Int -> Int
power n x
    | n == 0 = 1
    | even n = square (power (n `div` 2) x)
    | otherwise = x * power (n - 1) x
spower :: Int -> Code Int -> Code Int
spower n x
    | n == 0 = [|| 1 ||]
    | even n = [|| square $$(spower (n `div` 2) x) ||]
    | otherwise = [|| $$x * $$(spower (n - 1) x) ||]
```


## Staging generics-sop

## Basic idea

Structure is statically known, so rather than Sum (Product I) (Description a)
let us use
Sum (Product Code') (Description a)

## Staged conversions

```
class Generic a => SGeneric a where
sfrom ::
        Code a
    -> (Sum (Product Code') (Description a) -> Code r)
    -> Code r
sto ::
    Sum (Product Code') (Description a)
    -> Code a
```


## Example

The function sfrom introduces case analysis and passes the representation to the continuation:
instance SGeneric Animal where

```
sfrom x k =
    [|]
        case $$x of
            HoppingAnimal n d ->
            $$(k (Z (Code [|| n ||]:* Code [|| d ||]:* Nil)))
            WalkingAnimal n i ->
            $$(k (S (Z (Code [|| n ||]:* Code [|| i ||]:* Nil))))
        |]
sto x = ...
```


## Staged generic encode

```
gencode ::
    \forall a . (Generic a, All (All Serialise) (Description a))
    => a -> Encoding
gencode x =
    collapsesum
    (czipWith3sum @(All Top)
        (\ (K a) (K i) encs ->
            K ( encodeListLen (a + 1)
            <> encodeWord i
            <> mconcat (collapse Product encs)
                )
            )
            (constructorArities @a)
            (constructorNumbers @a)
        (cmapsop @Serialise (\ (I y) -> K (encode y)) (from x))
    )
```


## Staged generic encode

```
sgencode ::
    \forall a . (SGeneric a, All (All Serialise) (Description a))
    => Code (a -> Encoding)
sgencode =
    [|| \x -> $$(sfrom [|| x ||] $ \ x' ->
        collapse
            (czipWith3sum @(All Top)
            (\ (K a) (K i) encs -> let a' = a + 1 in
            K [|| encodeListLen a'
            <> encodeWord i
            <> $$(smconcat (collapse Product encs))
                II]
        )
        (constructorArities @a)
        (constructorNumbers @a)
        (cmapsop @Serialise
            (\ (Code y) -> K [|| encode $$y ||]) x')
        )
) II]
```


## Missing function

```
smconcat :: Monoid a => [Code a] -> Code a
smconcat [] = [|| mempty ||]
smconcat [x] = x
smconcat (x : xs) = [|| $$x <> $$(smconcat xs) ||]
```


## Status

## What (also) works

Implementing other staged generic functions:

- Deriving lenses (getters and setters).
- Generic equality and comparison.


## Current limitations

data IntList = IntCons Int IntList | IntNil
instance Serialise IntList where

## encode = \$\$(sgencode @IntList)

GHC stage restriction: instance for 'Serialise IntList 'is used in a top-level splice [...] and must be imported, not defined locally

## Current limitations

data IntList = IntCons Int IntList | IntNil
instance Serialise IntList where
encode = \$\$(sgencode @IntList)

GHC stage restriction: instance for 'Serialise IntList ' is used in a top-level splice [...] and must be imported, not defined locally
data Option a = None | Some a
instance Serialise (Option a) where
encode = \$\$(sgencode @(Option a))

No instance for (Serialise a) arising from a use of'sgencode '

## Open questions

Are the conversion functions sfrom and sto sufficient?
E.g. quadratic code size for generic equality.

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Are the conversion functions sfrom and sto sufficient?
E.g. quadratic code size for generic equality.

Can we transfer all the other known techniques from staging SYB (Yallop)?

